



# Guaranteeing Dependency Enforcement in Software Updates

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# Outline

- Motivating scenario
- Software update models
- The proposed solution
- More complex scenarios
  - Improved solutions and limitations

# Motivating Scenario

- A flock of drones is accomplishing a mission
- Some of the drones' missions needs to be reconfigured automatically by the Central/Regional Operations Office.
- The set of drones to be reconfigured depend on their configuration/status.
  - Availability of specific tools/weapons
  - Fuel level
  - Role in the current mission



# Underlying requirements

- On-the-fly mission reconfiguration
  - No human intervention/authorization
  - Mission specification is a sensitive information!
- Drones might be captured
  - Exposed material should not compromise future missions
- No or minimal hardware support for security

# Software Update scenario

- *Hardware components* are specified by means of *software* modules (drivers)
- *Missions* are specified by means of *software* modules
- *Drones' selection* is defined by the *installation policy*
  - *A mission can be installed if and only if the configuration meets the installation policy!*
- Updates might be:
  - Human assisted: Installation of new hardware component
  - Automatic: On-the-fly mission reconfiguration.

# Existing solutions

- Different software update **models**/systems are available:
  - Centralized *a-la* iOS
  - Semi-decentralized: Android
  - Fully distributed: Linux
- Different **tools** to secure updates:
  - Authentication: Package Signature, Code Signing
  - Confidentiality: Channel encryption
  - Dependency enforcement: optional!



# The Actors: Linux philosophy

- (Multiple) Distribution Server: Trusted component
  - Produces properly formatted software packages
- Mirror Server: Untrusted component
  - Used by the DS to distribute software packages
- Device: Partially trusted
  - If captured discloses key material
- The adversary
  - Controls the communication channel
  - Has full access to captured devices.



# Updicator [ABCSS14]

- Idea: Each *device* is identified by *static attributes*
  - O.S., CPU Type, Speed, etc.
- Uses **CP-ABE** to protect software **confidentiality**
  - Installation policy is embedded in the ciphertext
  - Each device holding the proper set of attributes can decrypt the software and install it.
- **Problem:** Attributes **updates** requires one of the following:
  - Existence of a **single** centralized key **authority**
  - Complete system **profiling** (transfer of all keys to an update authority)
  - **Interactive** protocol for key updates.





# Our requirements

- **Attributes** are *dynamic* in nature
  - Each new installed software enhances the device with a new attribute
- **Multiple** ‘key’ generation **authorities**
  - Each device might install software from different vendors
- **Non-interactive** key-updates
- **Enforce installation policy** in a “*strong sense*”
  - It is not possible to bypass it
- Typical security requirements
  - Software authenticity, integrity, confidentiality and freshness



# A hidden assumption

*If a package A depends on package B, the software vendor has already installed package B*

- Otherwise, how can the vendor test software A ?
- When the package is being assembled, the vendor knows the *attributes* of all the required packages.

# The idea

- Each software has two associated random keys
  - An **encryption key**, used to encrypt it
  - A **package key**, the ‘attribute key’
- Package creation:
  - Encrypt (software, **package key**) using the **encryption key**
  - Share **encryption key** using the installation policy as an access structure.
  - Encrypt **share** for ‘software j’ with the **package key** of software j.
- Installation is possible iff the set of installed packages satisfies the installation policy.



# The Package creation protocol

1. *Generate a random encryption key  $r$*
2.  $(s_1, \dots, s_m) \leftarrow \text{Distribute}(r, A_\phi)$
3. For  $j=1$  to  $m$  do
4.      $e_j \leftarrow \text{Encrypt}_{k_j}(s_j)$
5. *Generate a random package key  $k$*
6. Package  $p = (\text{name}, \text{timestamp}, \Delta, \text{metatada}, \text{package key } k, \text{software})$
7.  $E \leftarrow \langle \langle (n_1, e_1), \dots, (n_m, e_m) \rangle, A_\phi, \text{Encrypt}_r(p) \rangle$
8.  $M_E \leftarrow \text{CreateMetaData}(n, t, \Delta, M_s, E)$
9.  $\sigma_M \leftarrow \text{Sign}_{SKV}(M_E)$
10.  $\sigma \leftarrow \text{Sign}_{SKV}(E)$
11. Send  $(E, \sigma), (M_E, \sigma_M)$  to Mirror Servers



# Package Installation 1/2

Obtain  $(E, \sigma)$  and  $(M_E, \sigma_M)$  (possibly from different servers)

*/\* Authenticity, Integrity and Freshness Verification \*/*

1. if  $(\text{Verify}_{\text{pkV}}(E, \sigma) = \perp) \vee (\text{Verify}_{\text{pkV}}(M_E, \sigma_M) = \perp)$   
then Reject
2. if (E does not Match  $M_E$ )  $\vee$  ( $M_E$  Not Fresh)  
then Reject

*/\* Decryption Key Reconstruction \*/*

1. Parse E as  $\langle ((n_1, e_1), \dots, (n_m, e_m)), A_\varphi, \text{Encrypt}_k(p) \rangle$
2. for  $i = 1, \dots, m$  do  
if  $(n_i, k_i) \in \text{Installed Packages}$  then  $s_i = \text{Decrypt}_{k_i}(e_i)$   
else  $s_i = \perp$
3.  $r \leftarrow \text{Reconstruct}(A_\varphi, (s_1, \dots, s_n))$



# Package Installation 2/2

/\* Software Installation and Key Database update \*/

1.  $p \leftarrow \text{Decrypt}_r(E)$

2.  $M_s \leftarrow \text{ExtractMetadata}(M_E)$

if  $(p \neq \perp) \wedge (p \text{ matches } M_s)$  then

Parse  $p$  as  $(n,t,\Delta,M_s,k,s)$

Install  $s$

Add  $(n,k)$  to the set of Installed Packages

# Security Properties

- **Confidentiality**: guaranteed by package encryption
- **Integrity, Authenticity**: Enforced using standard tools
  - Hash functions, signatures, certificates...
- **Freshness**: Defined by the software vendor
- **Policy enforcement**: guaranteed by impossibility of decrypting the package without the knowledge of the proper keys.



# Advantages

- Each package brings its own attribute, the package key.
- Attributes' updates are executed non-interactively
  - “Save the package key”.
- Multiple-vendor updates are possible.



# More complex scenarios

- What if the installation policy is not monotone ?
  - Software policies contain the ‘conflict’ clause
  - Only monotone access structures are possible for secret sharing schemes!

	#packages	#packages with conflicts	Percentage
CentOs	8652	377	4,4%
RHEL 7.1	4432	299	6,7%
openSUSE 13.2	5334	242	4,5%
Fedora 21	2477	127	5,1%



# More complex scenarios

- What if installation policy itself is sensitive ?
  - Our first solution assumes the access structure to be transferred in clear to the device

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# Policy (Partial) Hiding Protocol

- The vendor *locally* anonymizes package names and policy
  - *Arbitrarily* maps each package name to an integer.
  - Describes installation policy in DNF over the set of anonymized package names



# Policy (Partial) Hiding Protocol

- At installation time
  - For each clause in the DNF formula tries to decrypt each share using the keys in the local DB
  - Gets the key when it succeeds in decrypting all the shares in a clause
    - Depends on the number of packages installed on the device!



# Policy (Partial) Hiding Protocol

- In theory anonymity comes to at a huge price: efficiency
  - A DNF formula might be exponentially longer than a compact representation
  - The device needs to blindly search in the proper key.
- In practice the impact might be affordable
  - DNF expansion is due to multiple version
  - On average the DNF formula is 25 time larger than ‘compact’ representation



# Conclusions

- Presented a system that allows the enforcement of installation policies during software installation/updates
  - Multiple independent vendors
  - Non-interactive key updates
  - Installation policies depend on static properties, e.g., installed packages
- Presented an extension for partially hiding policies
- Started evaluating the performance of the scheme



# Future Work

- Allow dynamic policies
  - Depending on ‘fuel level’, ‘current position’
- Consider non-monotone installation policies
- Reduce the impact of anonymization for the installation phase.